



Interference Management for Distributed Parallel Applications in Consolidated Clusters

Jaeung Han¹, Seungheun Jeon¹, Young-ri Choi², and Jaehyuk Huh¹

¹School of Computing, KAIST ²School of Electrical and Computer Engineering, UNIST

UC)is:C



- Resource contention causes performance interference
 - Last level caches, limited memory bandwidth, etc
- In single-node applications, the effect of intra-node interference is bounded within the system (node)



UC)is:C



- Resource contention causes performance interference
 - Last level caches, limited memory bandwidth, etc
- In single-node applications, the effect of intra-node interference is bounded within the system (node)



UC)is:C



- Resource contention causes performance interference
 - Last level caches, limited memory bandwidth, etc
- In single-node applications, the effect of intra-node interference is bounded within the system (node)



UCIIS T



- Resource contention causes performance interference
 - Last level caches, limited memory bandwidth, etc
- In single-node applications, the effect of intra-node interference is bounded within the system (node)
- In distributed applications, the interference effect from participating systems can interact with each other



Interference in Distributed Applications

 Execution time increases by interference in participating nodes vary by application characteristics



JCIST



Challenges in Distributed Applications

- Can we estimate performance impact of interference for distributed parallel applications?
- Two challenges
 - Interference in a subset of nodes: interference propagation



Different levels of interference: interference heterogeneity



 We propose a profiling-based interference estimation method

Quantifying Interference within a Node

- Bubble-Up [MICRO'11, Mars et al.]
 - Profiling-based interference model for single-node applications
 - Estimate the performance of co-located applications based on per-application interference profiles
- Per-application interference profile
 - Sensitivity profile: performance sensitivity to various levels of interference from the co-runner
 - Pressure score: interference level generated by the application







Bubble-Up Interference Profile

- Interference intensity is quantified to interference *pressure score*
- Bubble generates tunable amounts of interference pressure
- Reporter measures the pressure score (interference intensity generated by the application)







Bubble-Up Interference Profile

- Interference intensity is quantified to interference *pressure score*
- Bubble generates tunable amounts of interference pressure
- Reporter: measure the pressure score (interference intensity generated by the application)

Interference Profile for Distributed Applications

- 1) Pressure Score
- 2) Interference Propagation Profile
- 3) Heterogeneity Conversion Policy

Resource	

Bubble's pressure

UNIST



Propagation in Distributed Applications

- Interference on a subset of nodes can slow down the execution progress in non-interfering nodes
- Interference propagation profile
 - Execution time changes by the number of interfering nodes
 - Each node suffers from the same level of interference



Common Interference Propagation Patterns



High propagation

- One interfering node affects the exec. time significantly
- 104.milc, 126.lammps ...

Proportional propagation

- Exec. time increases proportionally
- 113.GemsFDTD ...

- Low propagation
 - Resilient to the interference
 - Kmeans(HADOOP), PageRank(SPARK) ...





Reducing Profiling Runs

- Binary-optimized
 - Shapes of curves are similar, regardless of pressure levels
 - Interpolate the exec. time from # of interfering nodes and pressure levels







Interference Propagation

- Binary-optimized
 - Shape of curves are similar, regardless of bubble pressures
 - Extrapolate the exec. time from # interfering nodes and bubble pressures







Interference Heterogeneity

• Each node can suffer with different interference intensity



- Too large space for profiling all possible heterogeneous interferences
 - 4 nodes + 9 interference levels : 495
 - 8 nodes + 9 interference levels : 12,870
 - 32 nodes + 9 interference levels : 76,904,685





Interference Heterogeneity Profile

- Interference Heterogeneity Profile
 - Convert heterogeneous interference to an equivalent hypothetical run with homogeneous interference







Conversion Policies

- 4 available conversion policies
 - N max
 - N+1 max
 - All max
 - Interpolate
- Evaluate all policies during profiling runs, and pick the best one for each application
- Use random sampling to reduce the number of profiling runs





Conversion Policies

- 4 available conversion policies
 - N max
 - Considers only the worst interfering nodes
 - N+1 max
 - All max
 - Interpolate







Convert Policies

- 4 available conversion policy
 - N max
 - N+1 max
 - Augments N max policy
 - The rest of interfering nodes are merged to the same worst pressure
 - All max
 - Interpolate







Convert Policies

- 4 available convert policies
 - N max
 - N+1 max
 - All max
 - The worst pressure propagates directly to all nodes
 - Interpolate







Convert Policies

- 4 available convert policies
 - N max
 - N+1 max
 - All max
 - Interpolate
 - Average interference from all nodes







Selecting Optimal Conversion Policy

Evaluate 4 policies for each application



- Select the best policy for each application
- Achieve less than 9% average error





Performance Estimation Steps

- Building interference profile for each application
 - 1. Build interference propagation profile (binary-optimized profiling)
 - 2. Measure interference intensity generated from the application (pressure score)
 - 3. Find the best heterogeneity conversion policy (random sampling)
- Estimating application execution time in a consolidated cluster
 - 1. For each node, find the interference intensity from the co-runner
 - 2. Apply the heterogeneity conversion policy, and find a hypothetical run with homogeneous interference
 - 3. Use the propagation profile to estimate the final execution time





Validation Results

- All possible pairwise combinations of workloads in consolidated runs
- The average error for each application against all the other applications as the co-runner
- Most of the workloads have less than 10% errors







Two Case Studies

- Placement for performance
 - Maximize the overall cluster throughput
 - Selected 10 workload combinations
 - Use simulated annealing(SA) as placement algorithm

- QoS-Aware placement
 - 1 target workload + 3 different co-runners
 - Provide QoS guarantee for the target workload
 - Compare to zero interference run
 - Use *SA* under the QoS Constraints as placement algorithm





Two Case Studies

- Placement for performance
 - Maximize the overall cluster throughput
 - Selected 10 workload combinations
 - Use simulated annealing(SA) as placement algorithm

- QoS-Aware placement
 - 1 target workload + 3 different co-runner
 - Provide QoS guarantee for the target workload
 - Compare to zero interference run
 - Use *SA* under the QoS Constraints as placement algorithm





Placement Results



- Best : the best placement based on performance estimation
- Random : Average result of 5 random placements
- Worst : the worst placement based on performance estimation





Results from Amazon EC2

• Validation for larger systems

Workload	Best Policy	Avg. error(%)	Std. dev.
M.milc	N+1 max	12.01	7.27
M.Gems	N+1 max	11.49	6.28
M.zeus	ALL max	6.40	4.52
M.lu	N max	5.28	4.36





Conclusion



- Proposed a profiling-based interference estimation for distributed applications
 - Extended the *Bubble-Up* technique
- Per-application interference profile
 - Pressure score + propagation profile + heterogeneity conversion
- Limitation 1: Static profiling
 - Assume a priori knowledge of each application
 - Cannot reflect dynamic changes
- Limitation 2: Pairwise interaction
 - Up-to two applications can be co-located on each node