

Fairness-oriented Scheduling Support for Multicore Systems

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Multicores have Uneven Capabilities



ΚΔΙST

Throughput-maximizing Scheduling

Max-perf: maximize the overall throughput

[Kumar'03], [Koufaty'10], [Kwon'11], [Saez'10], [Shelepov'09], [Craeynest'12]



KAI:

Fairness-Maximizing Scheduling

• Max-fair: maximize the fairness [Kwon'11]



- 1) Equal share of fast core to all apps
- 2) Equal share of slow core to all apps

KAIS

Max-Fair vs. Max-Perf

Fairness



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Throughput

Max-Fair vs. Max-Perf



KAIS1

Fairness-oriented Scheduling

Fairness

First, Fairness,

Then, Throughput

Policies

Min-fair: Guarantee MinFairness, then Throughput ↑ **Similar-min-fair:** Min-fair + Improve Uniformity

Implementation issues

Estimate accurate value of fast core speedup Schedule as core share

Min-Fair Scheduling

• Guarantee minimum performance as users want

KΛ

Similar-Min-Fair Scheduling

Similar speedup apps → same throughput ↑

- 1) Administrator sets *similarity* and *target* of MinFairness
- 2) Min-fair scheduling with *target*
- 3) Group similar apps (speedup difference < *similarity*)
- 4) Even fast core share out among apps in a group

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Sampling-based Speedup Estimation

- Force to run on both types of cores for each interval
 Similar to [Kumar '03]
- Measure performance using PMU (Performance Monitoring Units)
 - Performance metric: Instruction Per Seconds (IPS)

speedup = $\frac{\text{moving average of IPS on fast core}}{\text{moving average of IPS on slow core}}$

Architecture independent mechanism

Low error rate: avg. 2.70%, at worst 9.92%

Must run on both types → min-fair already does

- Incur frequent thread migration
- → Little impact on cores sharing a LLC [Craeynest '12]
- → Less than 1.8% at worst

Implementation on Linux Kernel

- Add fast_round and slow_round on thread context
 - +1 when run on fast/slow core for fast/slow_core_share*30ms
 - Progress together
 schedule as core share ratio
- Swap threads to balance fast_round and slow_round

- Measure IPS on each type of cores
- Core share update interval: 2 seconds

Evaluation Methodology

- Emulate uneven multicores using DVFS
 - AMD Phenom II: 2 fast cores and 4 slow cores

- **Big.LITTLE** architecture
 - Odroid-XU3 Lite: 4 big cores and 4 little cores
 - PMU not available → use offline speedup value
- Workloads: mix of SPECCPU2006
 - Run repeatedly until all apps finish at least once
 - Performance: execution time of the first run

Result: Diverse Applications

Guarantee MinFairness and then improve throughput

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Result: Similar Applications

Improve fairness without effect on throughput

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Result: All Mixes

Conclusion

- Fairness-oriented scheduling for uneven multicores

 First, Fairness, Then, Throughput
- Architecture independent speedup estimation
 High accuracy and little overhead
- Implemented on linux kernel 3.7.3
- Real machine results
 - MinFairness: mostly guaranteed (missed less than 3%)
 - Uniformity: avg. 24% ↑
 - Throughput: achieve 51% of Max-perf