

# Nested Enclave: Supporting Fine-grained Hierarchical Isolation with SGX

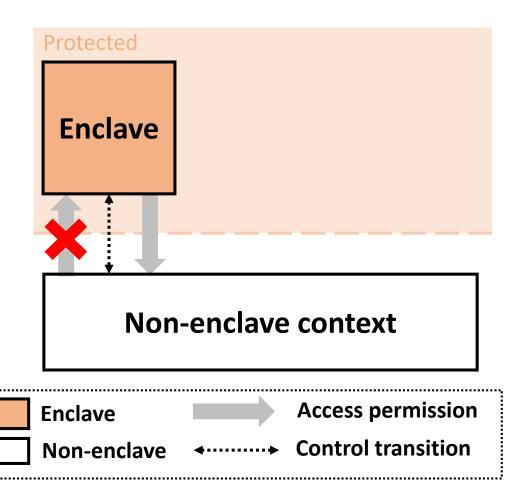
Joongun Park, Naegyong Kang, Taehoon Kim, Youngjin Kwon, Jaehyuk Huh

School of Computing, KAIST



#### SGX Trusted Execution Environment

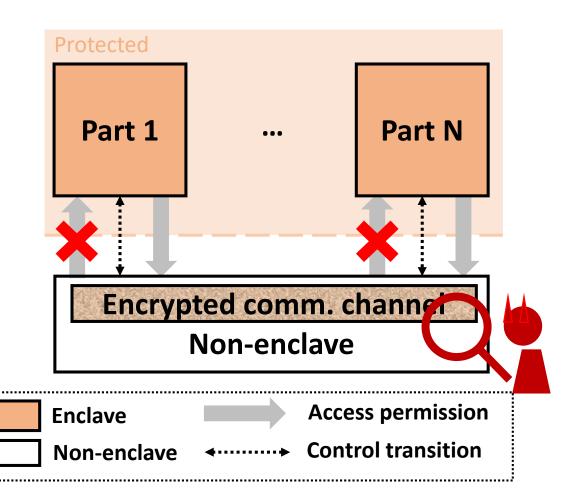
- Intel SGX
  - Provides trusted execution environments (TEE) called *Enclave*
  - Protected from malicious privileged SW
  - Guarantee confidentiality and integrity
  - Monolithic design





#### Need for Extension

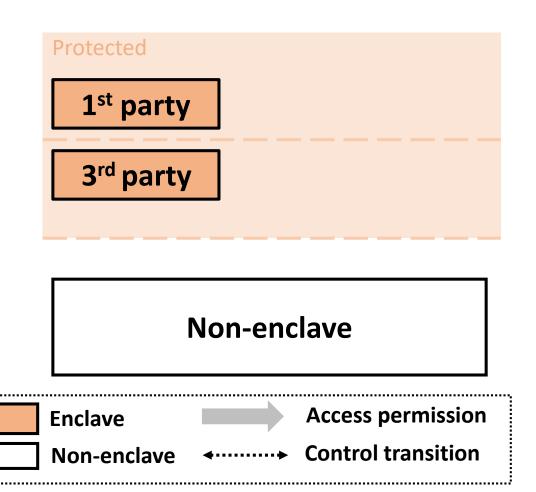
- Mutually distrustful parties [1]
  - •Multiple parties are involved for building an application
  - Use multiple enclaves
- Exposed communication channel [2]
  - Must protect against eavesdropping
  - Must detect silent drops





# Multi-level Security [3]

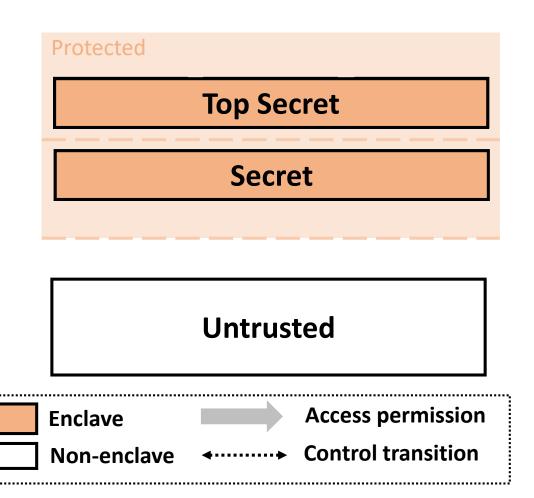
- Applications need multiple levels of security
  - Top secret / Secret
  - 1<sup>st</sup> party app / 3<sup>rd</sup> party library





# Compartmentalization

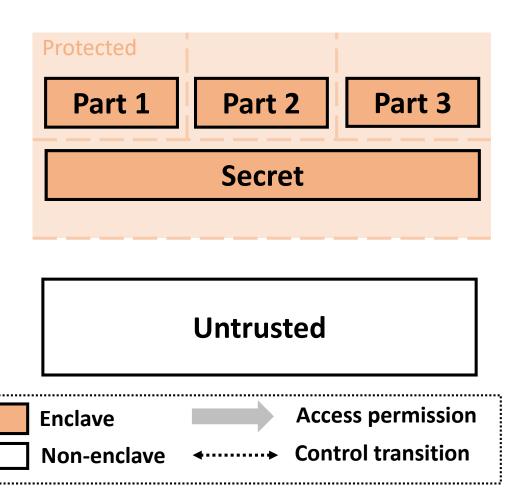
- Compartmentalization
  - Isolated peer compartments





# Compartmentalization

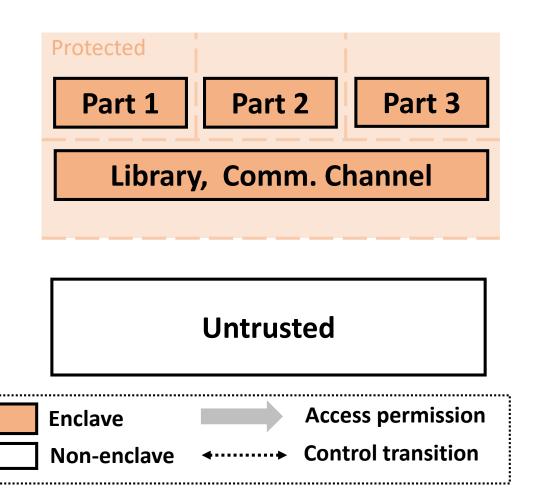
- Compartmentalization
  - Isolated peer compartments





# Compartmentalization

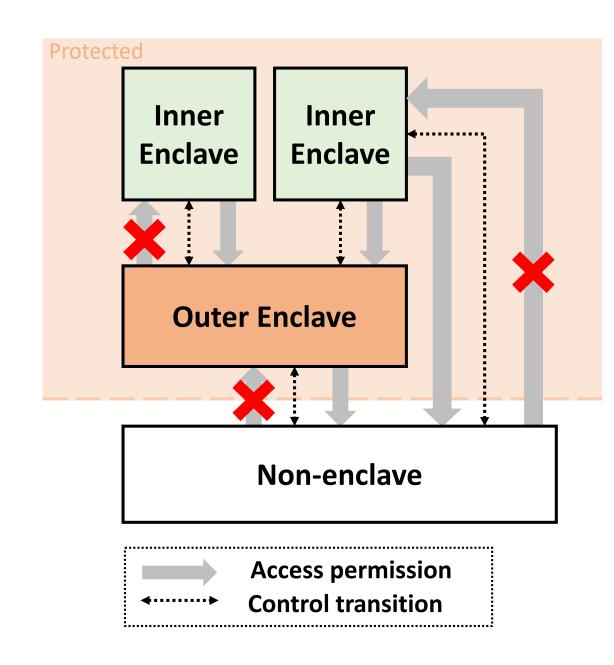
- Compartmentalization
  - Isolated peer compartments
- Sharing lower compartment
  - Shared library
  - Communication channel





#### Nested Enclave

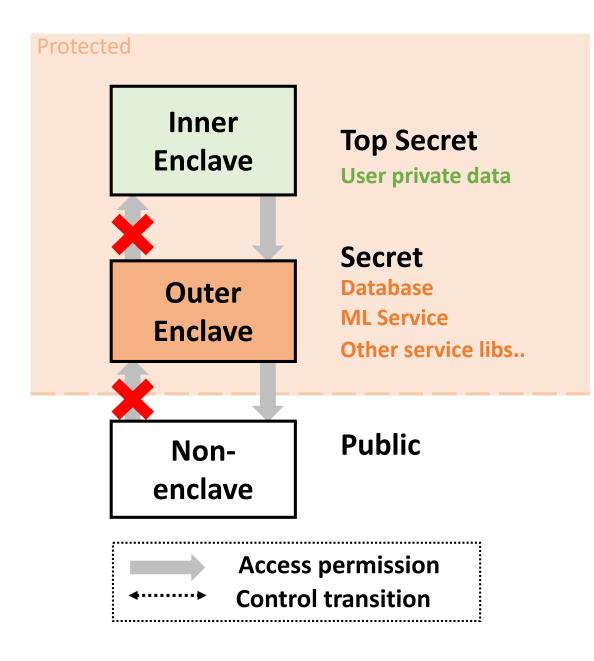
- Nested enclave is hardware extension to SGX
  - Inner Enclave
  - Outer Enclave





#### New semantics

- Hierarchical Isolation
  - Non-enclave context doesn't have access permission to both enclaves
  - Outer enclave doesn't have access permission to inner enclaves
  - Inner enclave has access permission to lower levels
- Supporting multi-level security

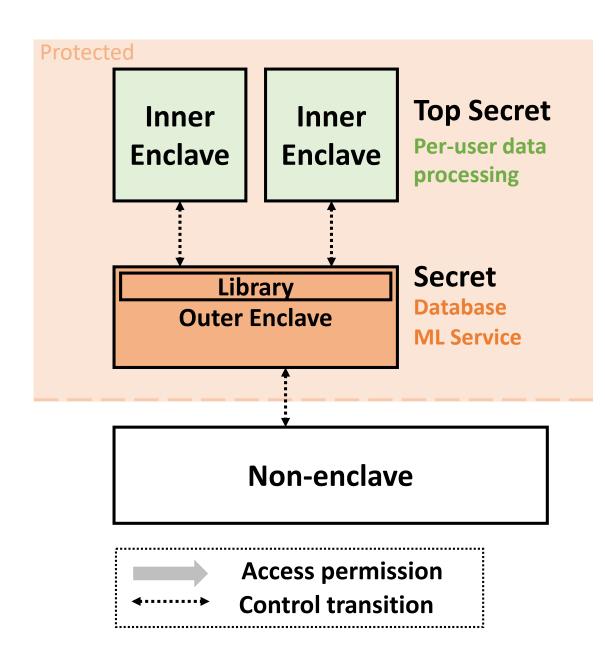




# New semantics (Cont'd)

- Compartmentalization
  - Isolation among Inner Enclaves

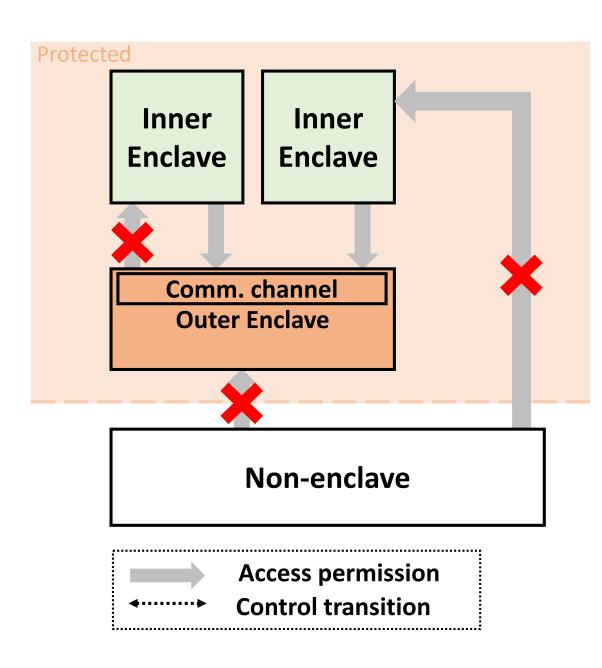
- Shared library with Outer Enclave
  - Reduced total memory usage
  - Inner enclave is protected from shared library





# New semantics (Cont'd)

- Secure communication channel in outer enclave
  - Hardware based protection via Memory Encryption Engine (MEE)
- Faster communication through caches
  - No encryption for data in caches
  - Enables faster data transfer [4]

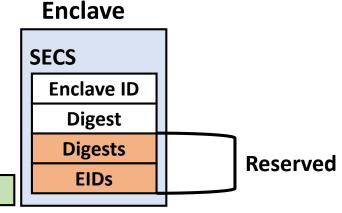




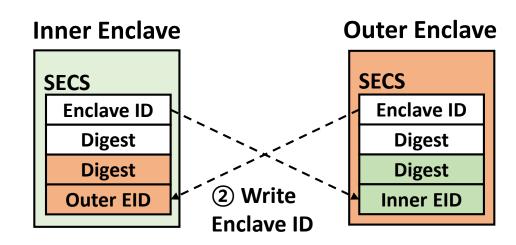
#### States and Association

- SGX Enclave Control Structure (SECS)
  - The state of an enclave
  - Use reserved fields to contain
    - Inner or outer enclave's digest
    - Inner or outer enclave's ID (EID)





- Association of inner-outer enclaves
  - Attest the pair enclave with digest
  - Write verified Enclave ID (EID) in SECS

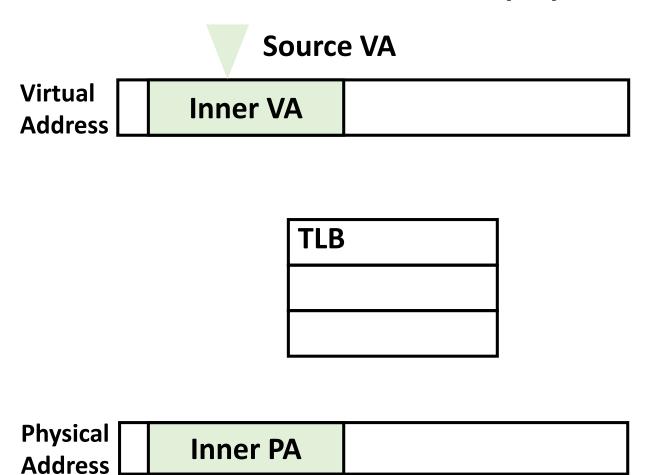




#### Memory Access Validation

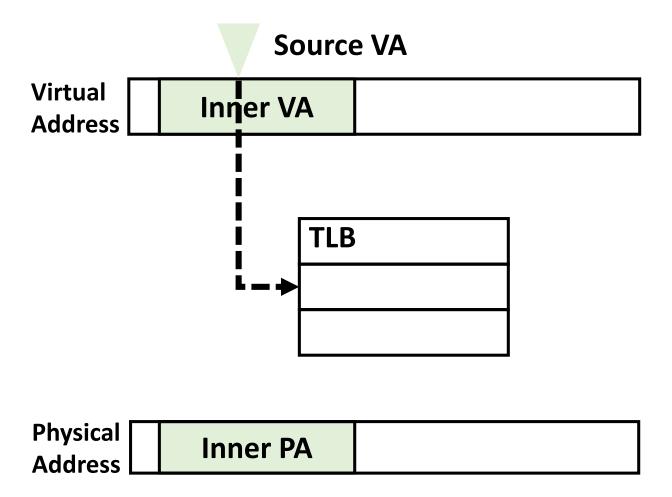
- SGX validates memory access during TLB miss handling
- EPCM (Enclave Page Cache Map) in SGX:
  - Meta data for each physical frame
  - Owner EID, VPN (virtual page number), etc.
- Invariant for security: TLB must contain only validated translation
- Cases
  - (A) Inner enclave accesses its enclave region
  - (B) Inner enclave accesses its outer enclave region
  - Others

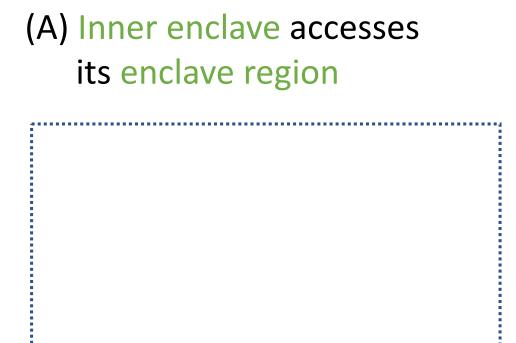




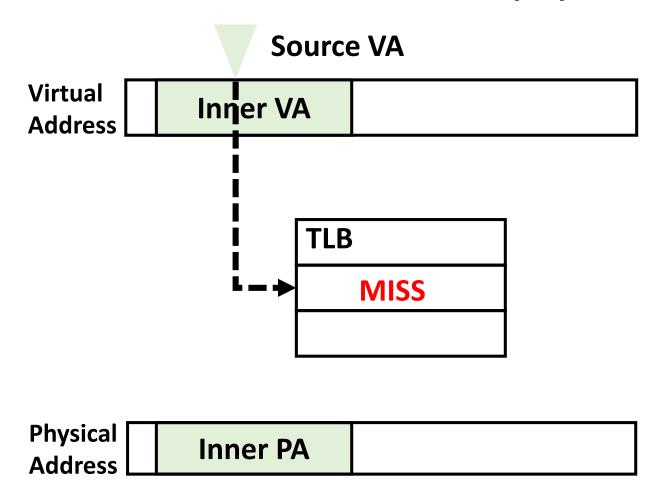
(A)	Inner encla	ave accesse	es
	its enclave	region	

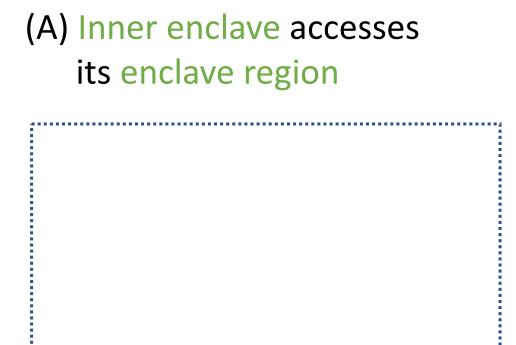




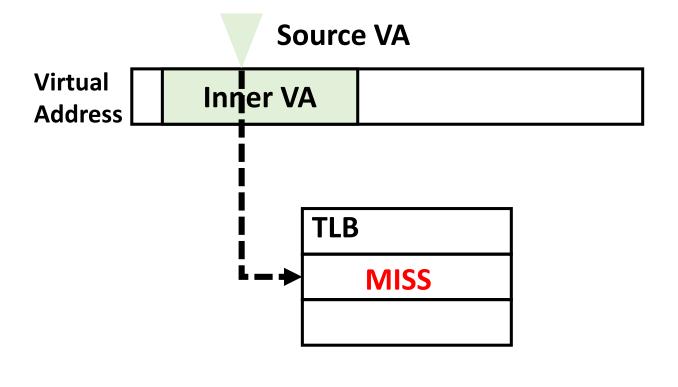








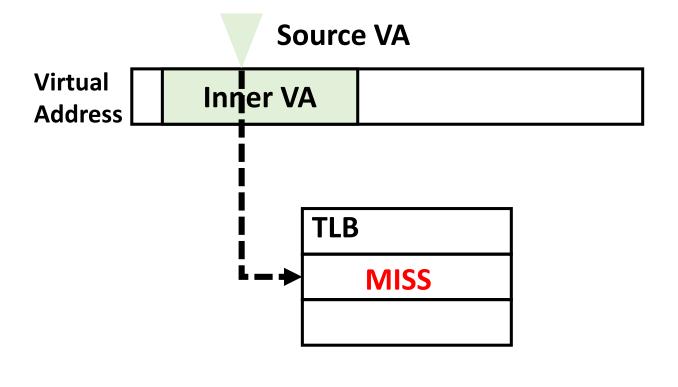






- (A) Inner enclave accesses its enclave region
  - (1) Target PA is Enclave PA?
  - (2) Check EPCM
    - Correct VPN?
    - (Owner EID == current EID)
      || (Owner EID == Outer EID)

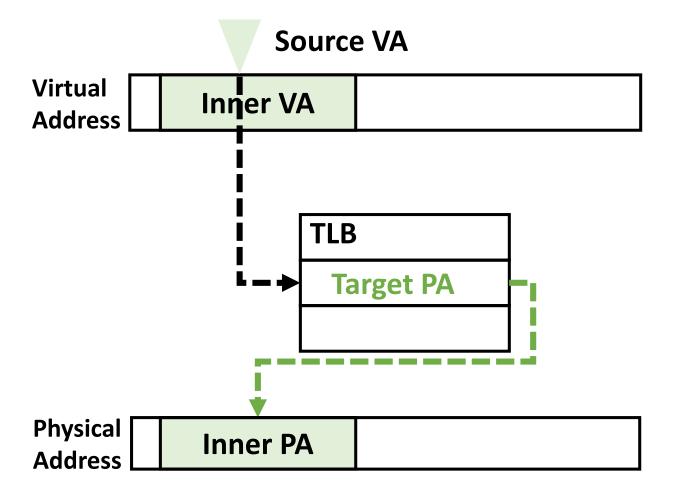






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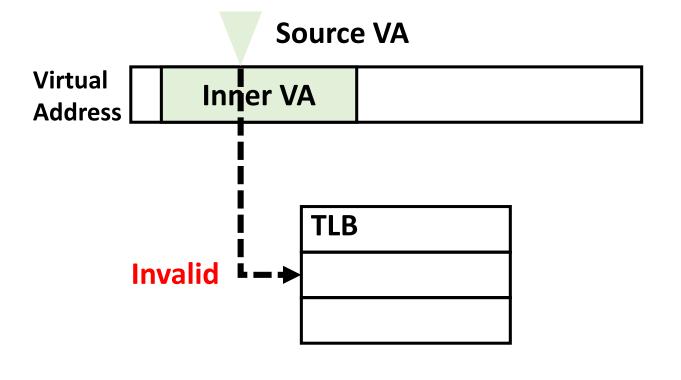


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YES => Insert TLB entry







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  - (1) Target PA is Enclave PA?
  - (2) Check EPCM
    - Correct VPN?

```
- (Owner EID == current EID)
|| (Owner EID == Outer EID)
```

YES => Insert TLB entry

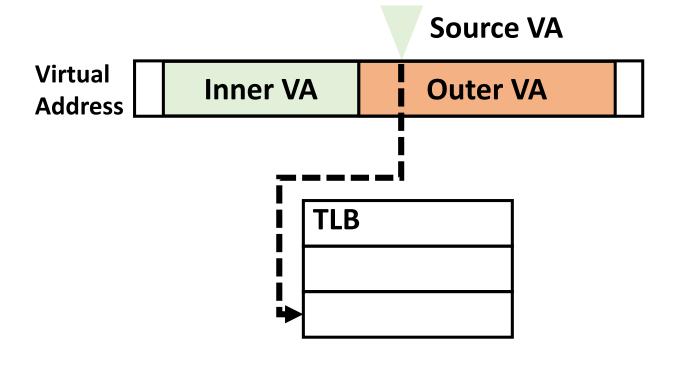
No => Invalid



Virtual Address	Inner VA	Outer VA	
	TLB		
Physical Address	Inner PA	Outer PA	

(B)	Inner enclave accesses		
	its outer enclave region		

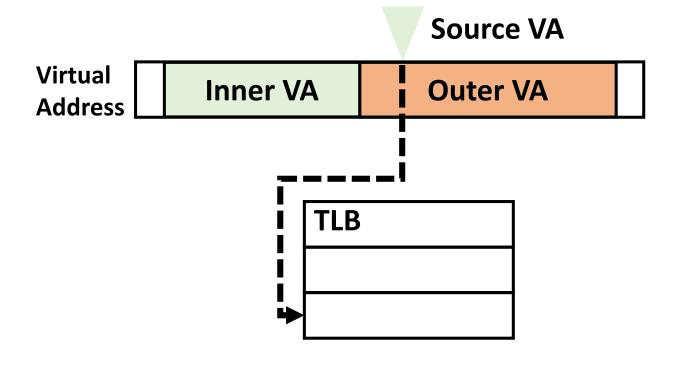




Physical Address Inner PA Outer PA

- (B) Inner enclave accesses its outer enclave region
  - (1) Target PA is Enclave PA?
  - (2) Check EPCM
    - Correct VPN?
    - (Owner EID == current EID)
      - || (Owner EID == Outer EID)

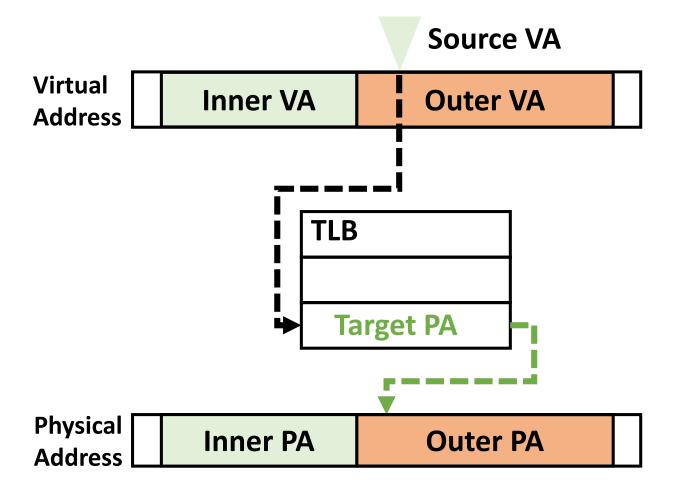




Physical Address Inner PA Outer PA

- (B) Inner enclave accesses its outer enclave region
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    - Correct VPN?
    - (Owner EID == current EID)
      - | | (Owner EID == Outer EID)



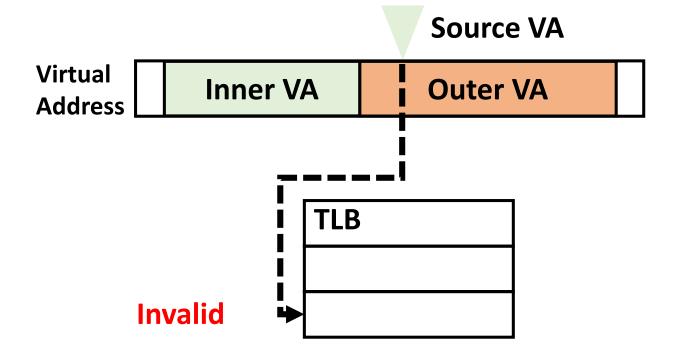


(B) Inner enclave accesses its outer enclave region

- (1) Target PA is Enclave PA?
- (2) Check EPCM
  - Correct VPN?
  - (Owner EID == current EID)
    - | | (Owner EID == *Outer EID*)

**YES => Insert TLB entry** 







(B) Inner enclave accesses its outer enclave region

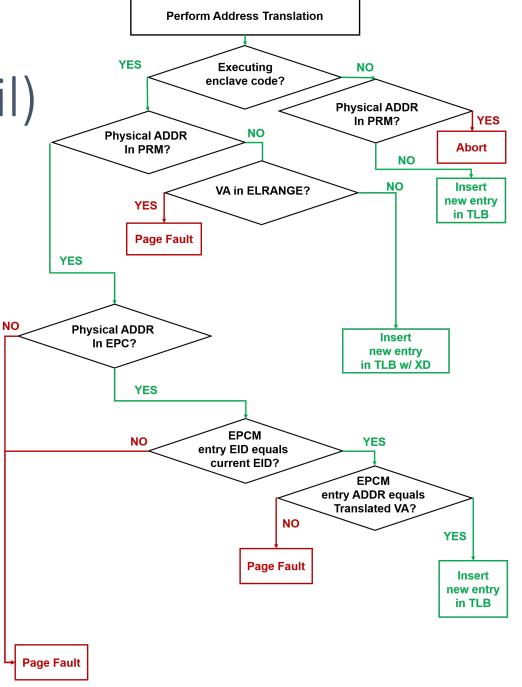
- (1) Target PA is Enclave PA?
- (2) Check EPCM
  - Correct VPN?
  - (Owner EID == current EID)

YES => Insert TLB entry

No => Invalid



Access Validation (detail)

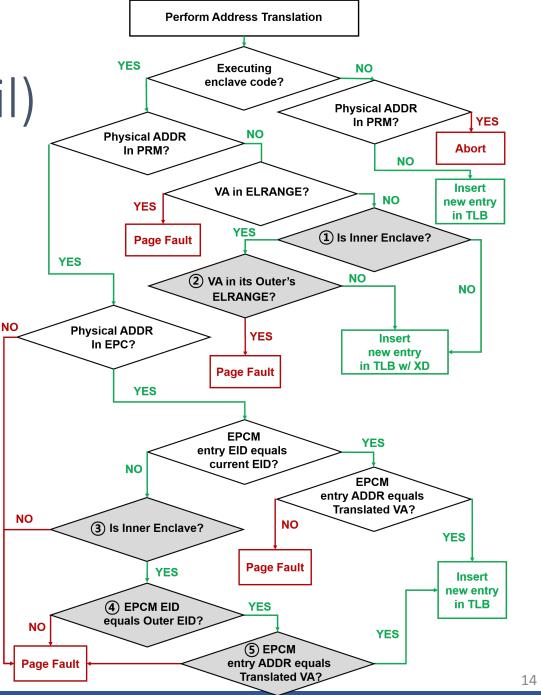


[4] Intel SGX Explained.

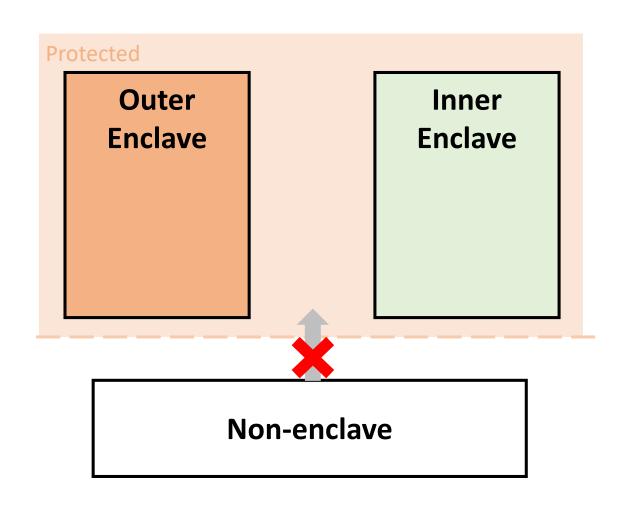


Access Validation (detail)

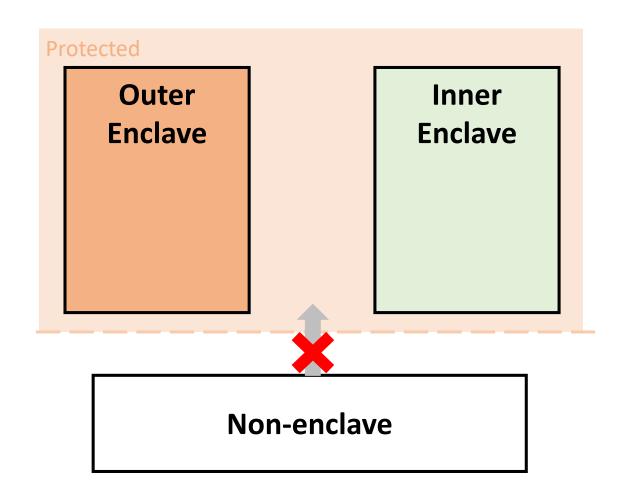
Modifications are marked in grey





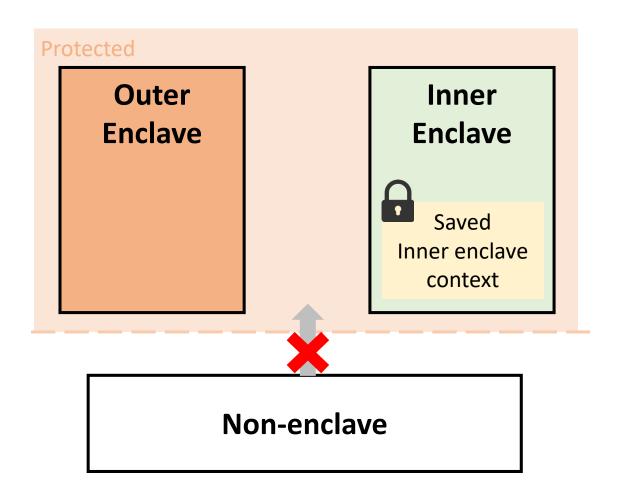






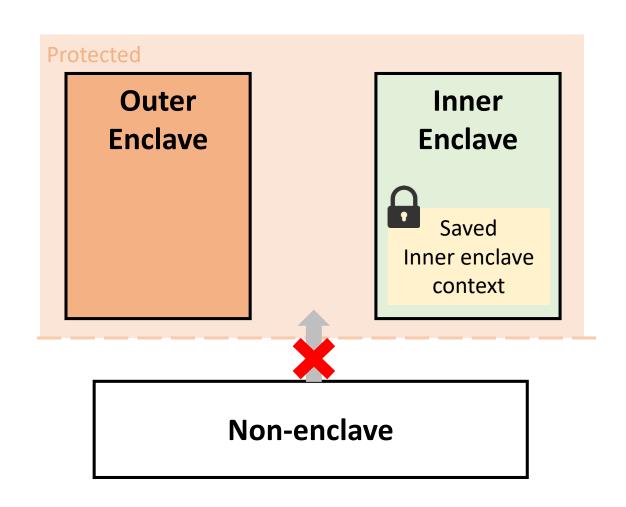


- Transition between Inner and outer enclaves
  - Save running context



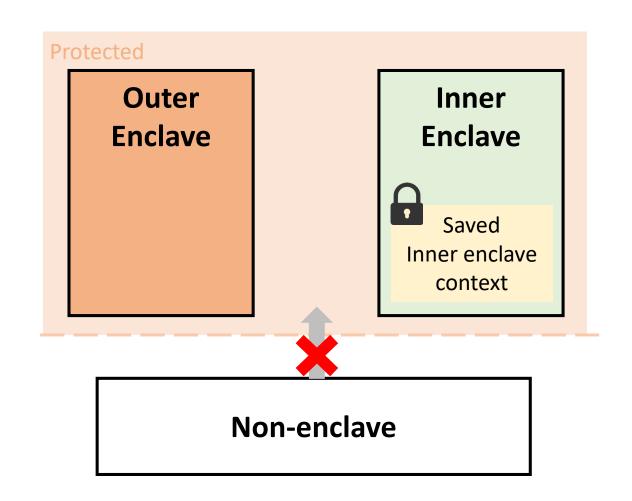


- Transition between Inner and outer enclaves
  - Save running context
  - Flush flags, register, and TLB



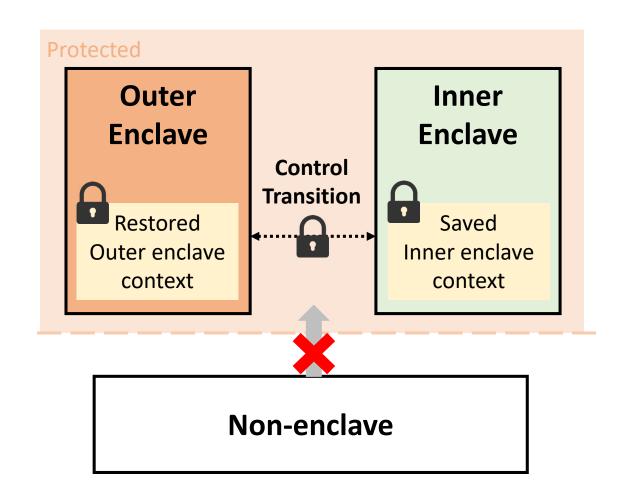


- Transition between Inner and outer enclaves
  - Save running context
  - Flush flags, register, and TLB
  - Check & sanitize parameters





- Transition between Inner and outer enclaves
  - Save running context
  - Flush flags, register, and TLB
  - Check & sanitize parameters
  - Restore target context if exists





# Evaluation Methodology

#### **Implementation**

- New instructions to SDK emulation
  - NEENTER, NEEXIT, NASSO, NEREPORT
- APIs
  - Nested ecall/ocall
  - Association

#### **Application porting**

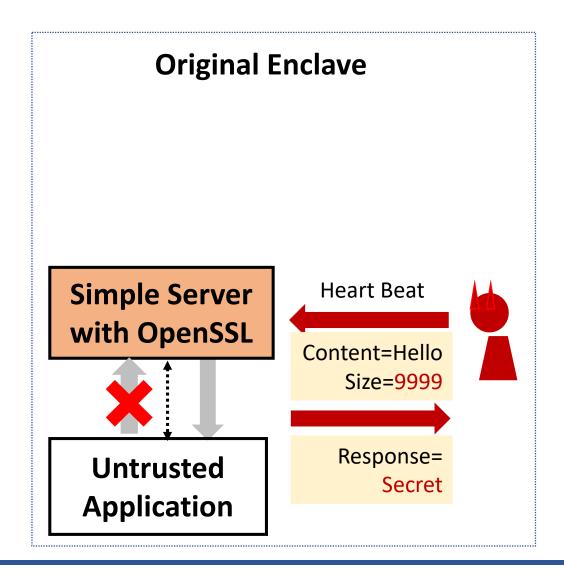
- Echo server with OpenSSL
- Query server with SQLite
- LibSVM (training, prediction)

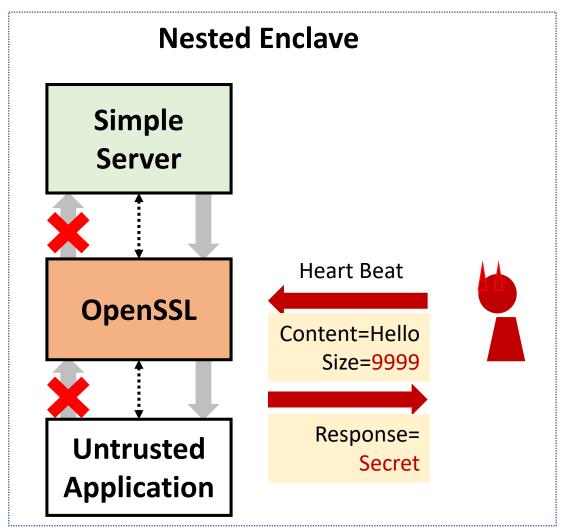
#### **Evaluation Environment**

- Intel i7-7700 64 GB DRAM
- Ubuntu 16.04, Linux kernel 4.13.0.
- Intel SGX SDK / driver v1.9



#### CASE 1: Heartbleed Attack







#### CASE 1. Hoarthlood Attack

G:INFO=YOUCANNOT

READITBECAUSEIFR

47 3B 49 4E 46 4F 3D 59 4F 55 43 41 4E 4E 4F 54

52 45 41 44 49 54 42 45 43 41 55 53 45 49 46 52

```
void info leak()
                                                                                                                                               In Simple Server
   char *secret;
   int size = 0x8000;
   secret = (char *) malloc (size);
far (int in 0) (int in 1) * 00100
      strlcpy (secret + i * 0x100, "ID=admin; PASSWORD=admin secure password 1337; EX=deadbeef; TOPSECRET=THISISACONFIDENTIALSTRING; INFO=YOUCANNOTREADITBECAUSEIFREEITAFTERUSETHIS; ", 0x100)
   free(secret);
                                                     READITBECAUSEIFR
Response
            52 45 41 44 49 54 42 45 43 41 55 53 45 49 46 52
                                                                                          Response
                                                                                           ....ID=admin:PA
                                                      SSWORD=admin sec
                                                      ure password 133
             75 72 65 5F 70 61 73 73   77 6F 72 64 5F 31 33 33
             37 3B 45 58 3D 64 65 61 64 62 65 65 66 3B 54 4F
                                                      7;EX=deadbeef;T0
                                                      PSECRET=THISISAC
                                                      ONFIDENTIALSTRIN
                                                                                          G; INFO=YOUCANNOT
                                                      READITBECAUSEIFR
                                                      EEITAFTERUSETHIS
                                                                                          00 49 44 3D 61 64 6D 69 6E 3B 50 41
                                                      .....ID=admin;P/
                                                      SSWORD=admin sec
            53 53 57 4F 52 44 3D 61 64 6D 69 6E 5F 73 65 63
            75 72 65 5F 70 61 73 73 77 6F 72 64 5F 31 33 33
                                                      ure password 133
             37 3B 45 58 3D 64 65 61 64 62 65 65 66 3B 54 4F
                                                      7;EX=deadbeef;T0
                                                      PSECRET=THISISAG
             50 53 45 43 52 45 54 3D  54 48 49 53 49 53 41 43
                                                                                           4F 4E 46 49 44 45 4E 54  49 41 4C 53 54 52 49 4E
                                                      ONFIDENTIALSTRIN
```



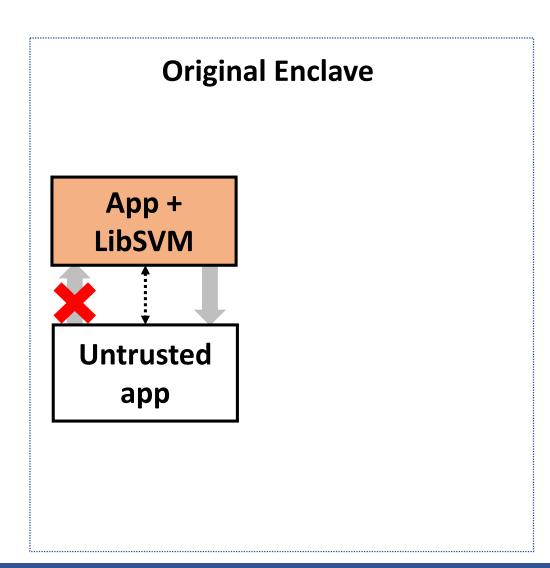
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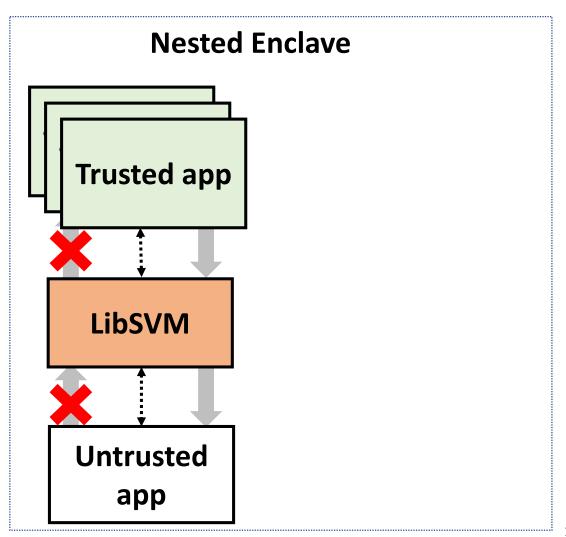
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   free(secret);
                                                               ure password 13
                                                               7;EX=deadbeef;T0
                                                               PSECRET=THISISAG
```

### Nested enclave protects the critical routine from the heartbleed attack

47 3B 49 4E 46 4F 3D 59 4F 55 43 41 4E 4E 4F 54 | G;INFO=YOUCANNOT 52 45 41 44 49 54 42 45 43 41 55 53 45 49 46 52 | READITBECAUSEIFR

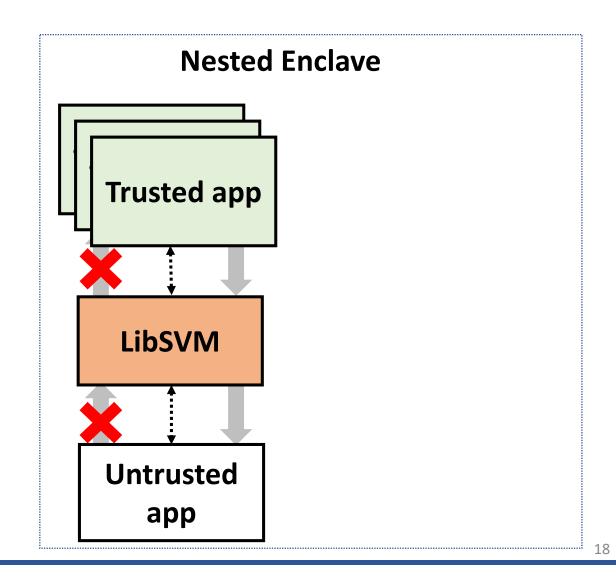






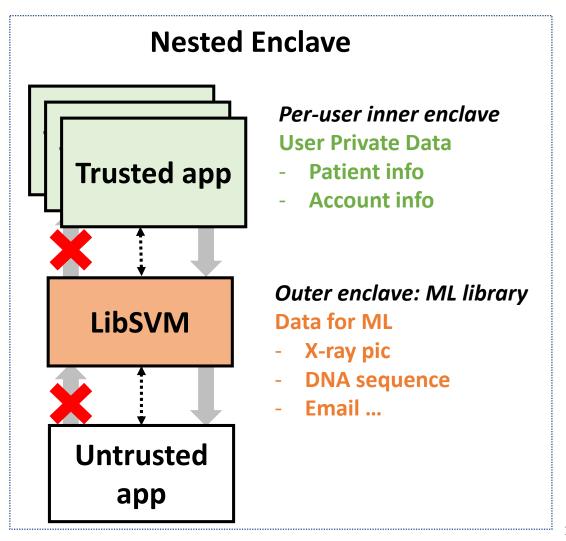


### **Original Enclave** Single enclave: Private processing+ App + ML library **LibSVM User Private Data** Patient info **Account info Untrusted** Data for ML - X-ray pic app **DNA** sequence Email ...



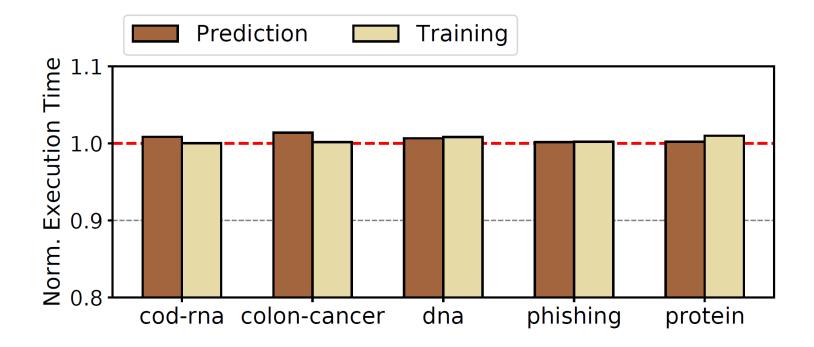


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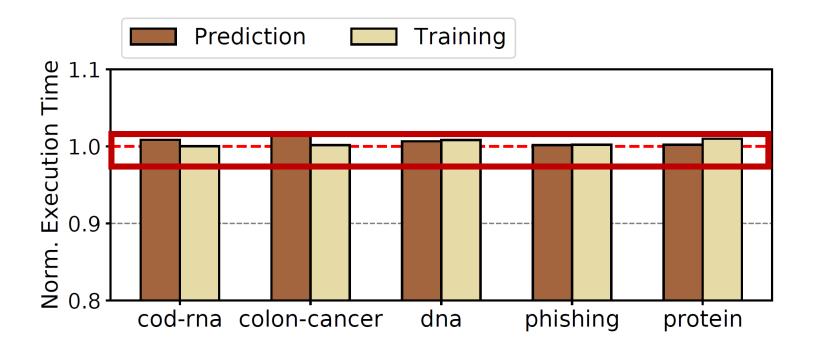


- Normalized execution time to the original enclave
- Overhead is relatively small ( < 2%)</li>





- Normalized execution time to the original enclave
- Overhead is relatively small ( < 2%)</li>





### More case studies in paper...

- Library sharing and time to loading enclaves
- Shared SQLite server
- Communication with intra- vs inter enclave channels



## Porting effort

Name	Modification	Modified LOC	Original
Echo server	C/C++ code	34	883
	EDL	10	28
	SGX-OpenSSL	0	507k
SVM-predict	C/C++ code	27	208
	EDL	10	49
	SGX-LibSVM	0	152k
SVM-train	C/C++ code	24	333
	EDL	10	41
	SGX-LibSVM	0	152k



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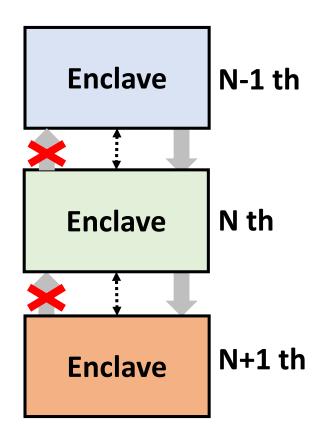


## Extending Nested Enclave



## Extending Nested Enclave

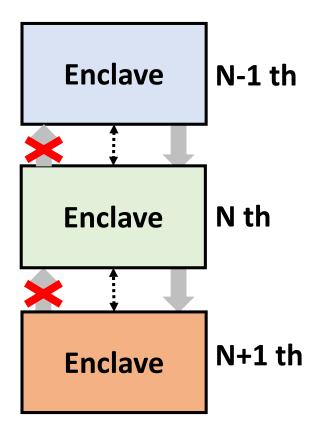
#### N-depth nesting



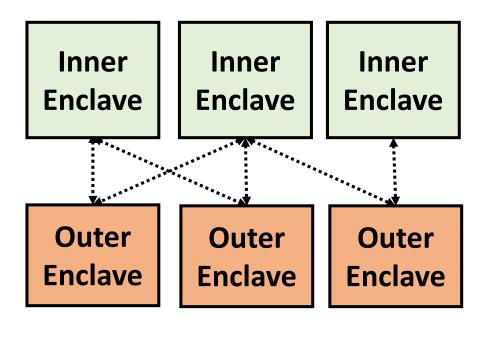


### **Extending Nested Enclave**

N-depth nesting



N to M relation





### Conclusion

- A new extension to the trusted execution environments for supporting multi-level security within TEE.
- The required hardware and software extensions for the nested enclave are small.
- Case studies on the nested enclave emulation framework.



# Thank you